

# Last time

- NAT
  
- Application layer
  - ◆ Intro
  
  - ◆ Web / HTTP

# This time

- Finish HTTP
- FTP

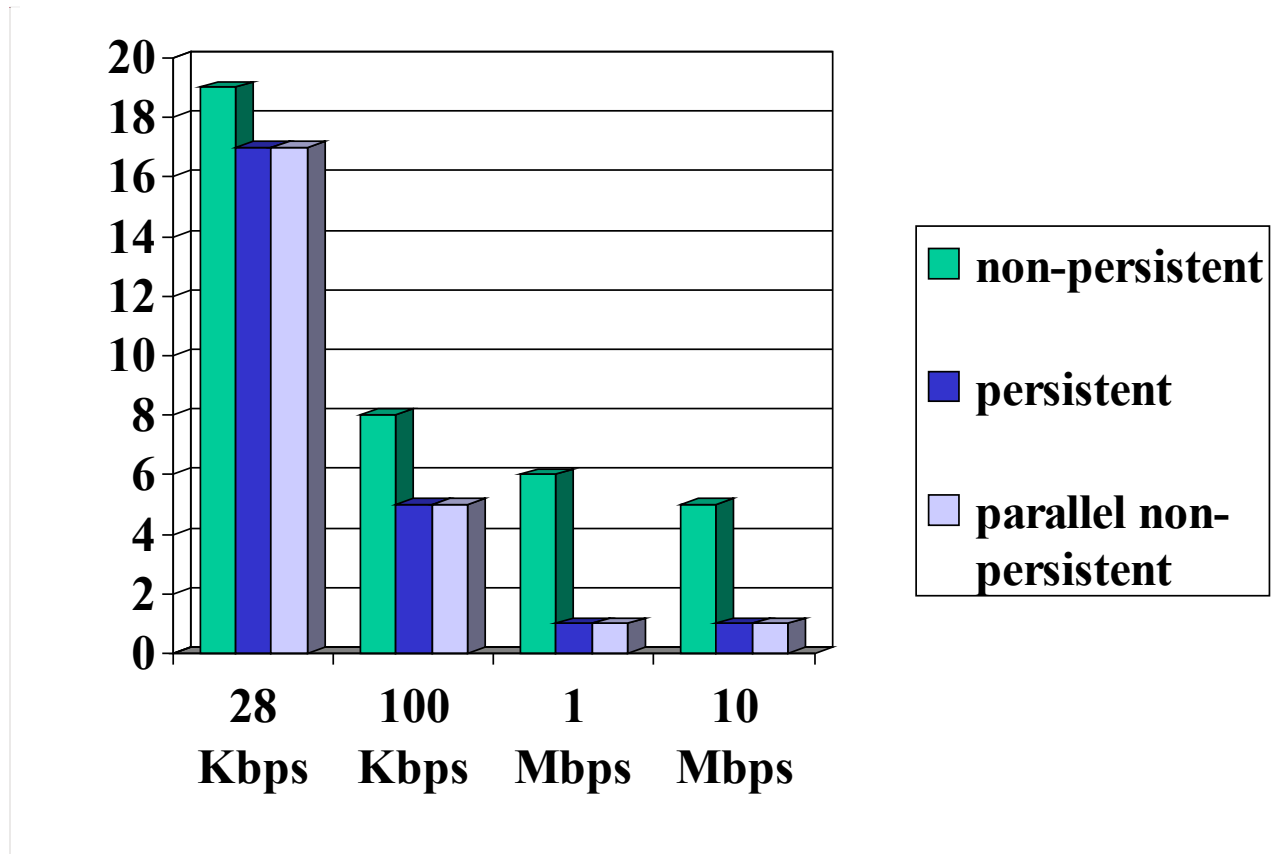
# HTTP Modeling

- Assume Web page consists of:
  - ◆ 1 base HTML page (of size  $O$  bits)
  - ◆  $M$  images (each of size  $O$  bits)
- Non-persistent HTTP:
  - ◆  $M+1$  TCP connections in series
  - ◆ *Response time =  $(M+1)O/R + (M+1)2RTT + \text{sum of idle times}$*
- Persistent HTTP:
  - ◆ 2  $RTT$  to request and receive base HTML file
  - ◆ 1  $RTT$  to request and receive  $M$  images
  - ◆ *Response time =  $(M+1)O/R + 3RTT + \text{sum of idle times}$*
- Non-persistent HTTP with  $X$  parallel connections
  - ◆ Suppose  $M/X$  integer.
  - ◆ 1 TCP connection for base file
  - ◆  $M/X$  sets of parallel connections for images.
  - ◆ *Response time =  $(M+1)O/R + (M/X + 1)2RTT + \text{sum of idle times}$*

See the applet on UW-ACE!

# HTTP Response time (in seconds)

RTT = 100 msec, O = 5 Kbytes, M=10 and X=5

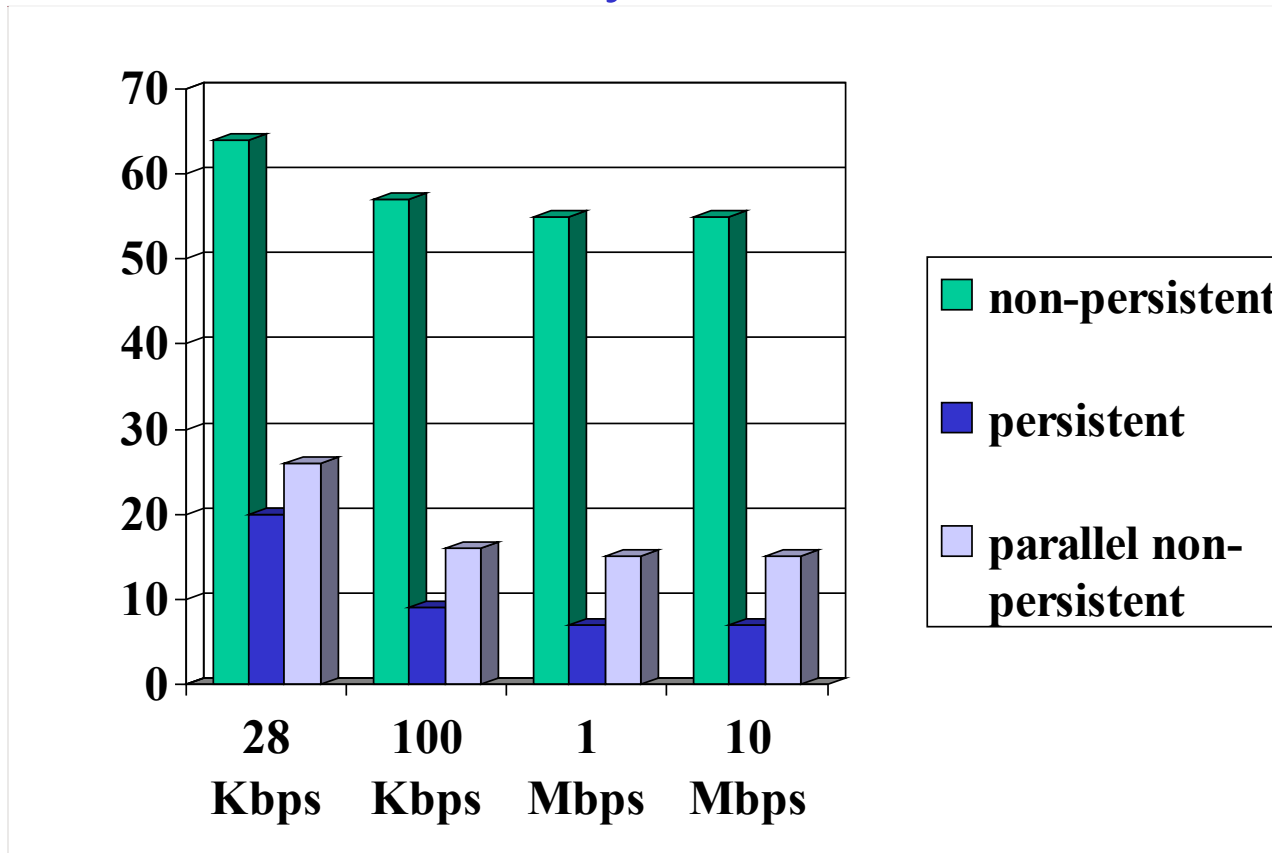


For low bandwidth, connection & response time dominated by transmission time.

Persistent connections only give minor improvement over parallel connections.

# HTTP Response time (in seconds)

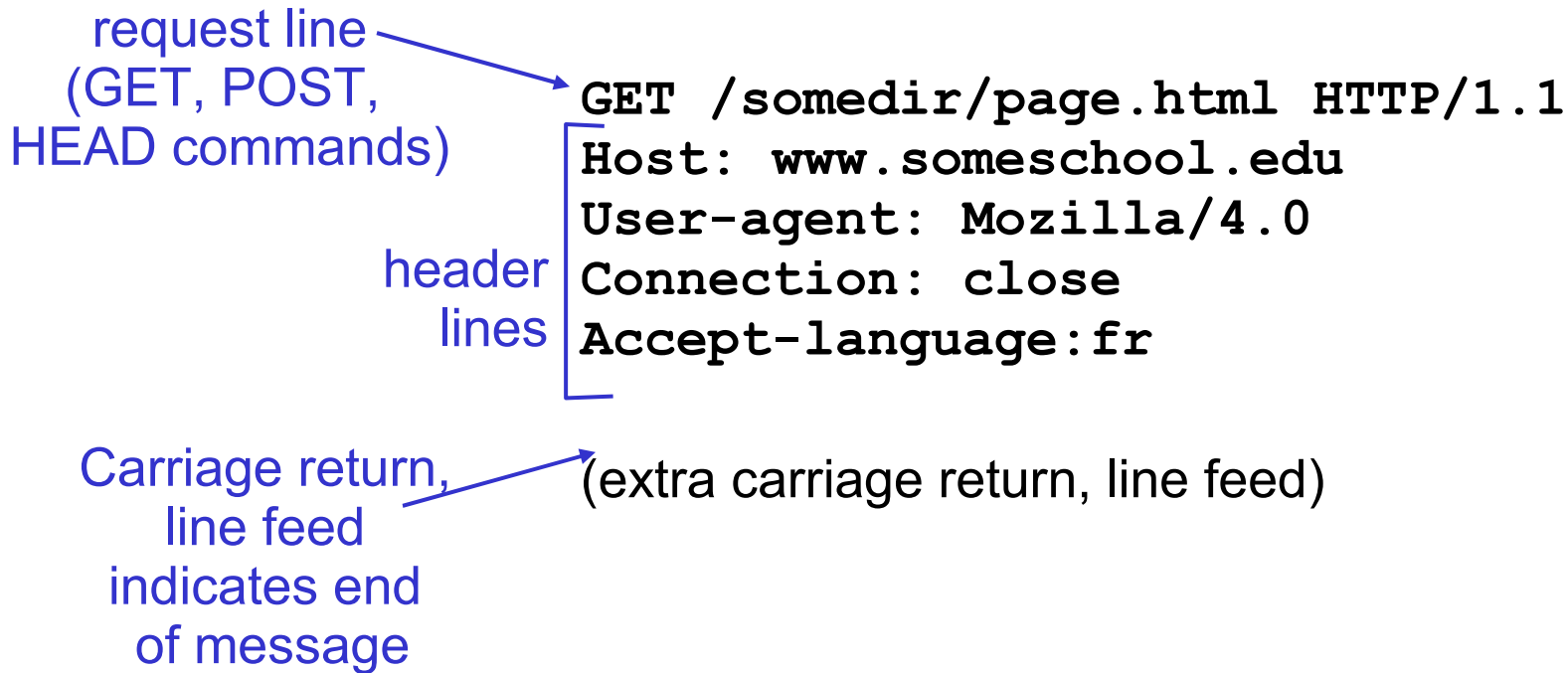
RTT = 1 sec, O = 5 Kbytes, M=10 and X=5



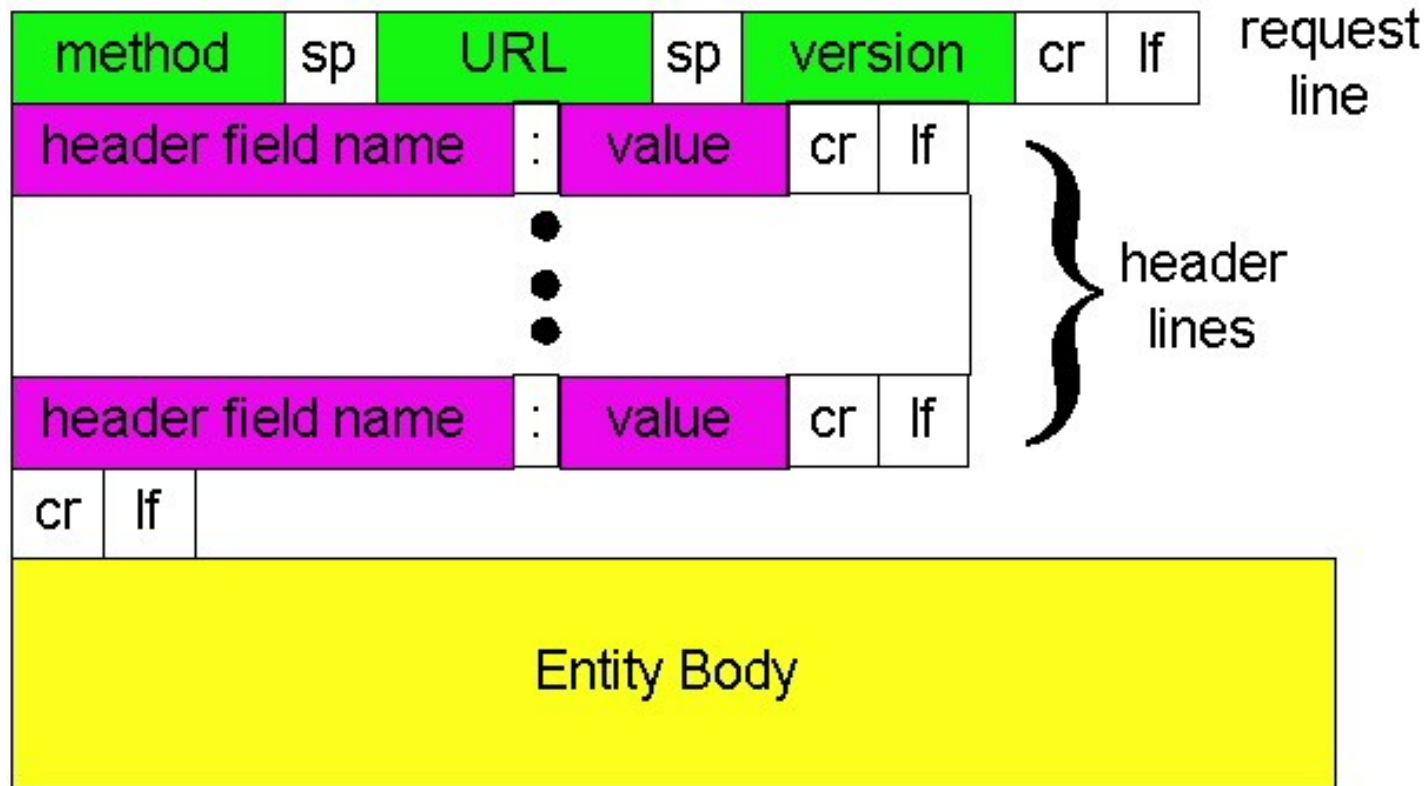
For larger RTT, response time dominated by TCP establishment & slow start delays. Persistent connections now give important improvement: particularly in high delay•bandwidth networks.

# HTTP request message

- Two types of HTTP messages: *request, response*
- **HTTP request message:**
  - ◆ ASCII (human-readable format)



# HTTP request message: general format



# Uploading form input

## POST method:

- Web page often includes form input
- Input is uploaded to server in entity body

## URL method:

- Uses GET method
- Input is uploaded in URL field of request line:

`http://www.somesite.com/animalsearch?monkeys&banana`



# Method types

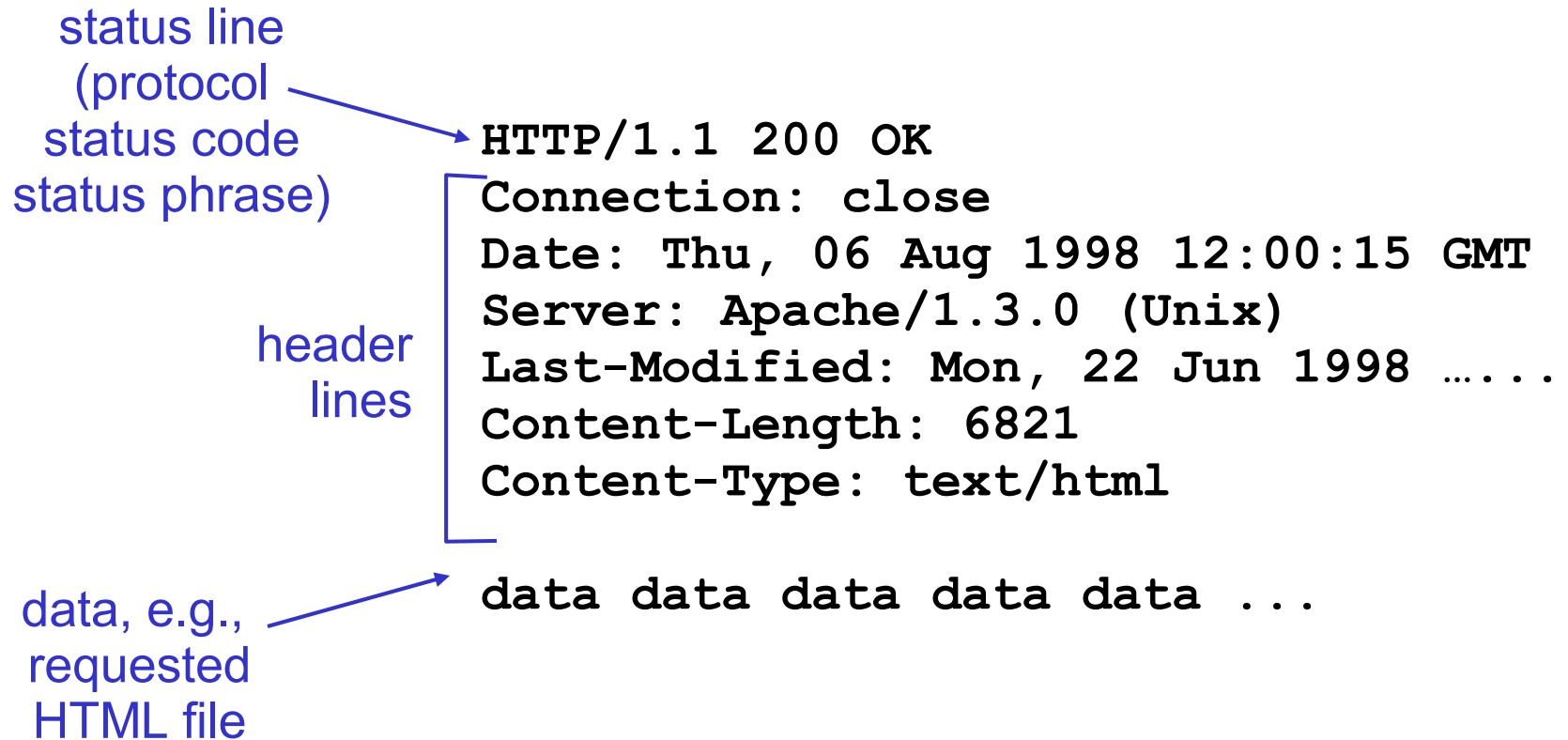
## HTTP/1.0

- GET
- POST
- HEAD
  - ◆ asks server to leave requested object out of response

## HTTP/1.1

- GET, POST, HEAD
- PUT
  - ◆ uploads file in entity body to path specified in URL field
- DELETE
  - ◆ deletes file specified in the URL field

# HTTP response message



# HTTP response status codes

In first line in server->client response message.

A few sample codes:

## **200 OK**

- ◆ request succeeded, requested object later in this message

## **301 Moved Permanently**

- ◆ requested object moved, new location specified later in this message (Location:)

## **400 Bad Request**

- ◆ request message not understood by server

## **404 Not Found**

- ◆ requested document not found on this server

## **505 HTTP Version Not Supported**

# Trying out HTTP (client side) for yourself

1. Telnet to your favorite Web server:

```
telnet cis.poly.edu 80
```

Opens TCP connection to port 80 (default HTTP server port) at cis.poly.edu. Anything typed in sent to port 80 at cis.poly.edu

2. Type in a GET HTTP request:

```
GET /~ross/ HTTP/1.1  
Host: cis.poly.edu
```

By typing this in (hit carriage return twice), you send this minimal (but complete) GET request to HTTP server

3. Look at response message sent by HTTP server!

# User-server state: cookies

Many major Web sites use cookies

## Four components:

- 1) cookie header line of HTTP *response* message
- 2) cookie header line in HTTP *request* message
- 3) cookie file kept on user's host, managed by user's browser
- 4) back-end database at Web site

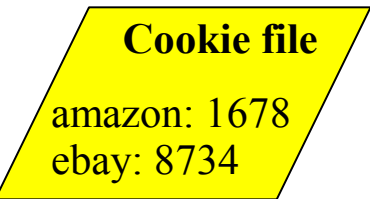
## Example:

- ◆ Susan access Internet always from same PC
- ◆ She visits a specific e-commerce site for first time
- ◆ When initial HTTP requests arrives at site, site creates a unique ID and creates an entry in backend database for ID

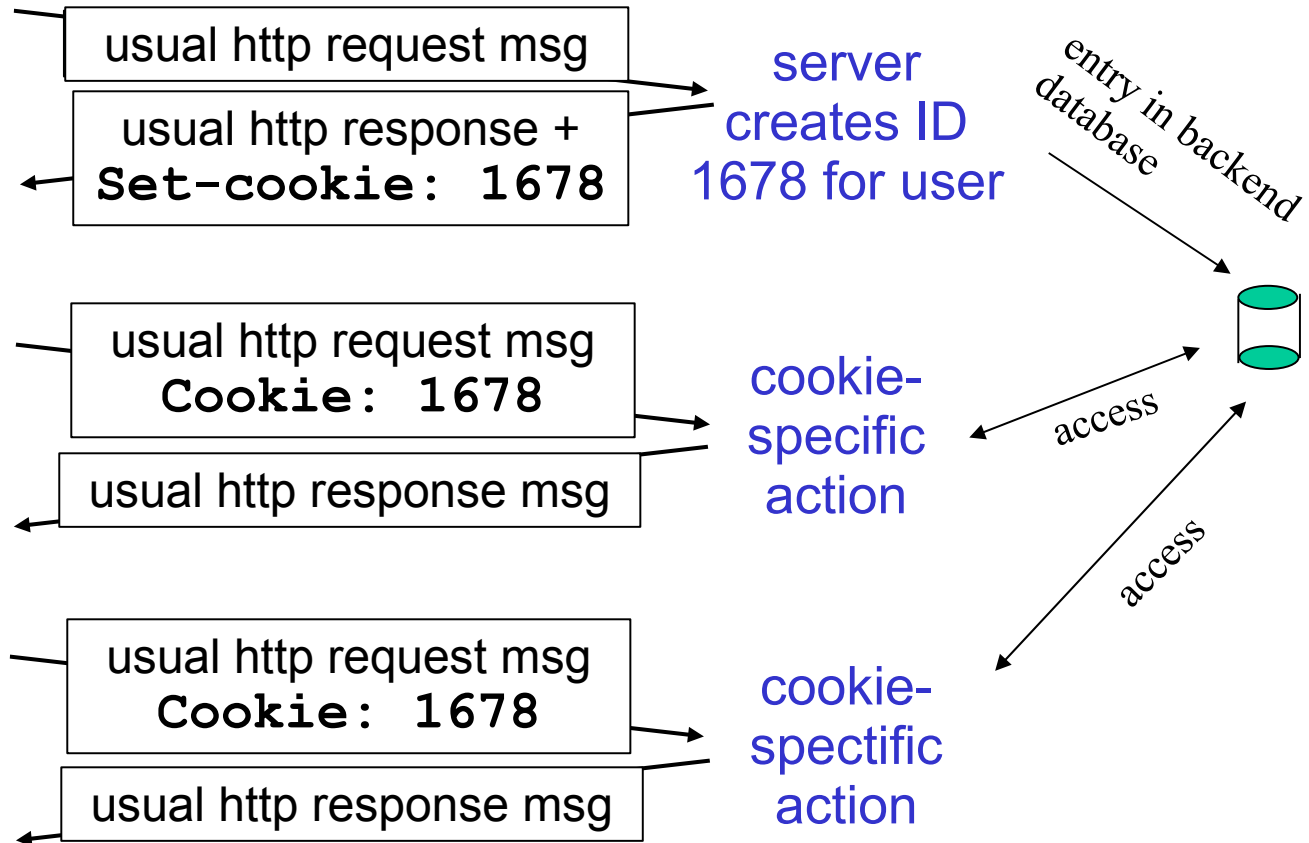
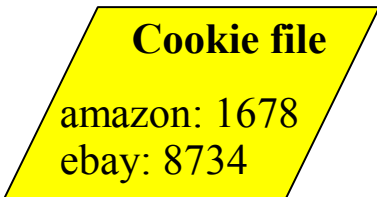
# Cookies: keeping "state" (cont.)

client

server



one week later:



# Cookies (continued)

## What cookies can bring:

- authorization
- shopping carts
- recommendations
- user session state (Web e-mail)

## How to keep “state”:

- Protocol endpoints: maintain state at sender/receiver over multiple transactions
- cookies: http messages carry state

aside

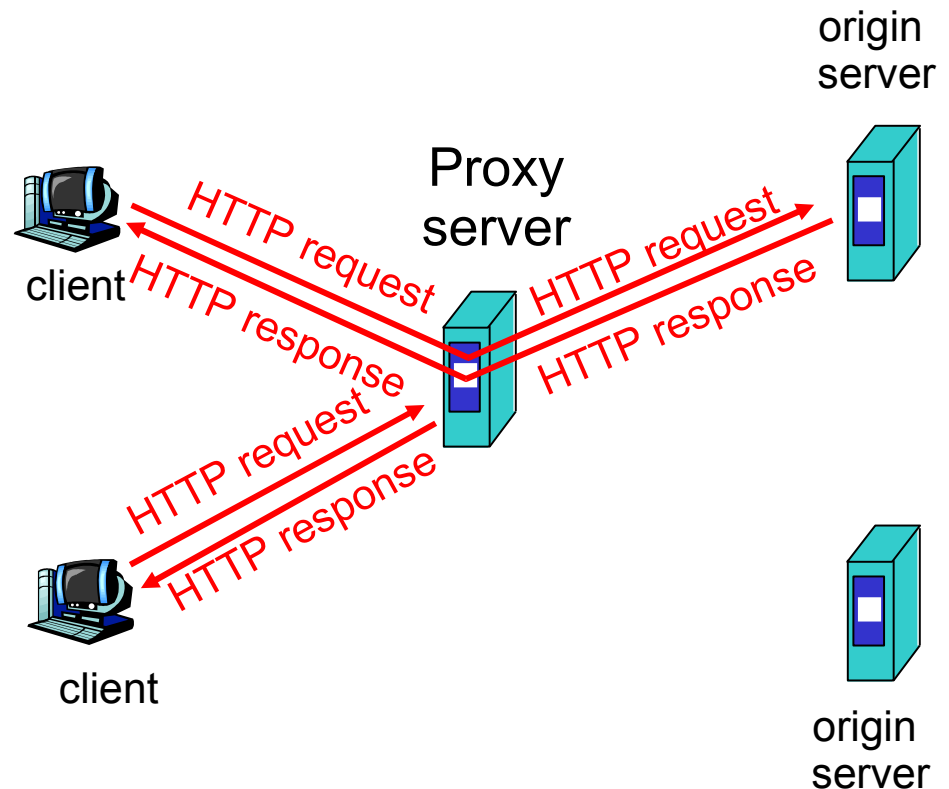
## Cookies and privacy:

- cookies permit sites to learn a lot about you
- you may supply name and e-mail to sites

# Web caches (proxy server)

**Goal:** satisfy client request without involving origin server

- user sets browser: Web accesses via cache
- browser sends all HTTP requests to cache
  - ◆ object in cache: cache returns object
  - ◆ else cache requests object from origin server, then returns object to client





# More about Web caching

- Cache acts as both client and server
- Typically cache is installed by ISP (university, company, residential ISP)

## Why Web caching?

- Reduce response time for client request.
- Reduce traffic on an institution's access link.
- Internet dense with caches: enables “poor” content providers to effectively deliver content (but so does P2P file sharing)

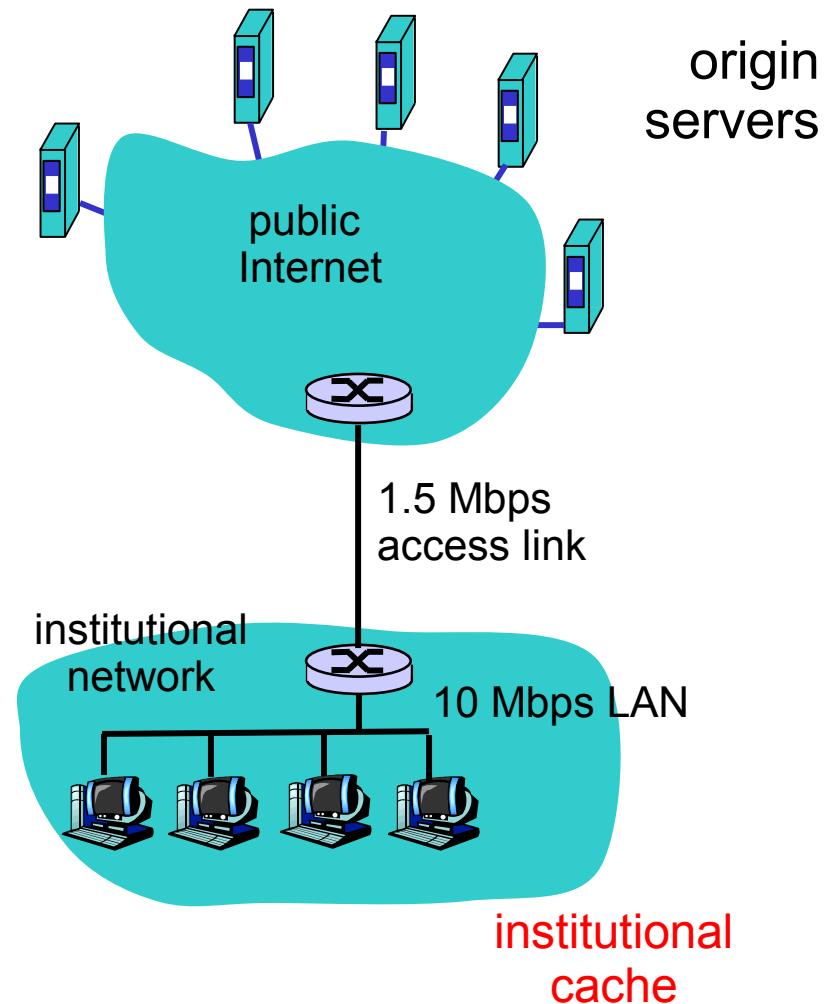
# Caching example

## Assumptions

- average object size = 100,000 bits
- avg. request rate from institution's browsers to origin servers = 15/sec
- delay from institutional router to any origin server and back to router = 2 sec

## Consequences

- utilization on LAN = 15%
- utilization on access link = 100%
- total delay = Internet delay + access delay + LAN delay  
= 2 sec + minutes + milliseconds



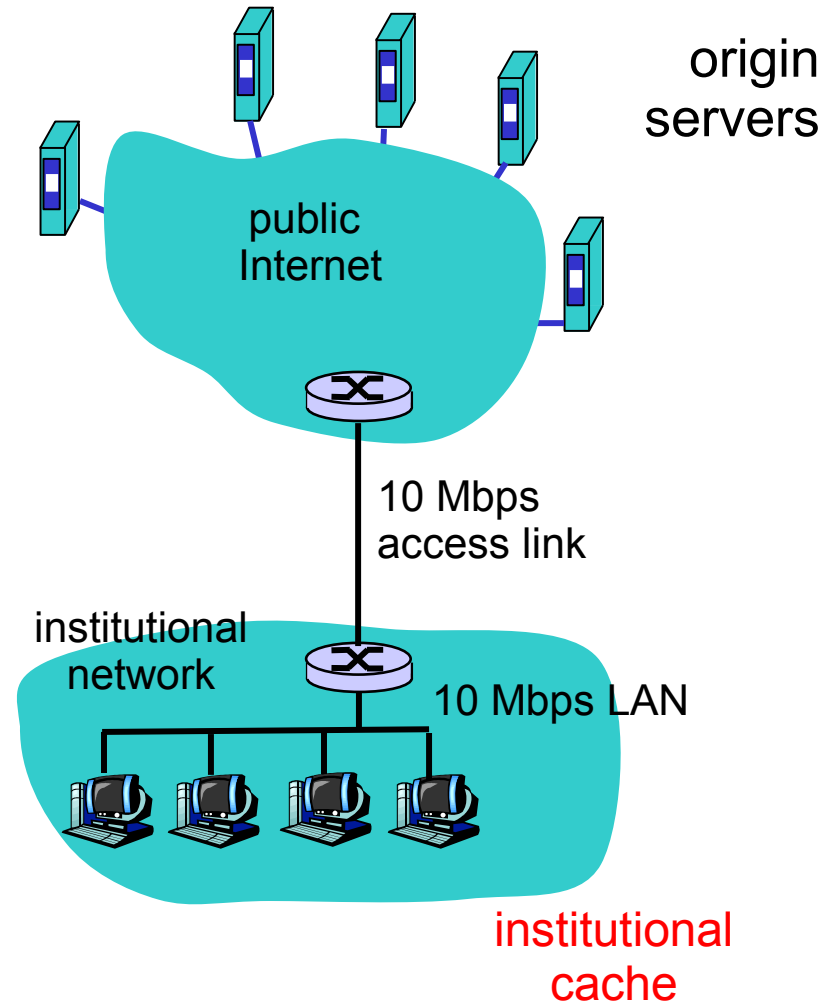
# Caching example (cont)

## Possible solution

- increase bandwidth of access link to, say, 10 Mbps

## Consequences

- utilization on LAN = 15%
- utilization on access link = 15%
- Total delay = Internet delay + access delay + LAN delay  
= 2 sec + msec + msec
- often a costly upgrade



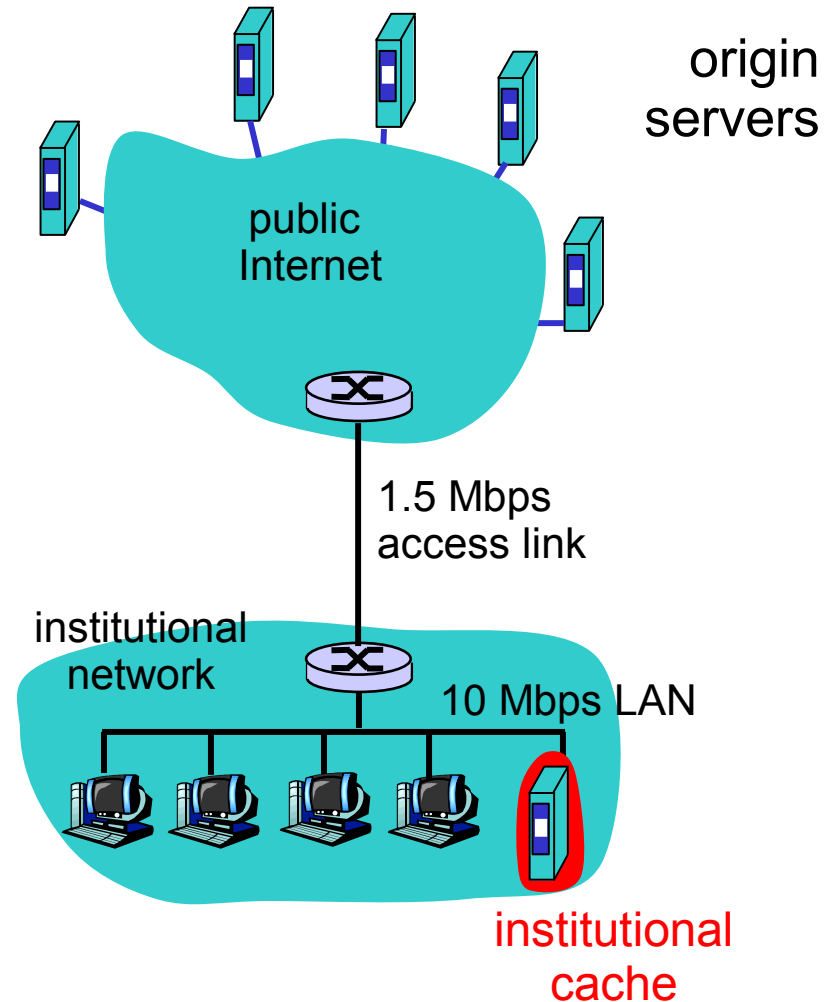
# Caching example (cont)

## Install cache

- suppose hit rate is .4

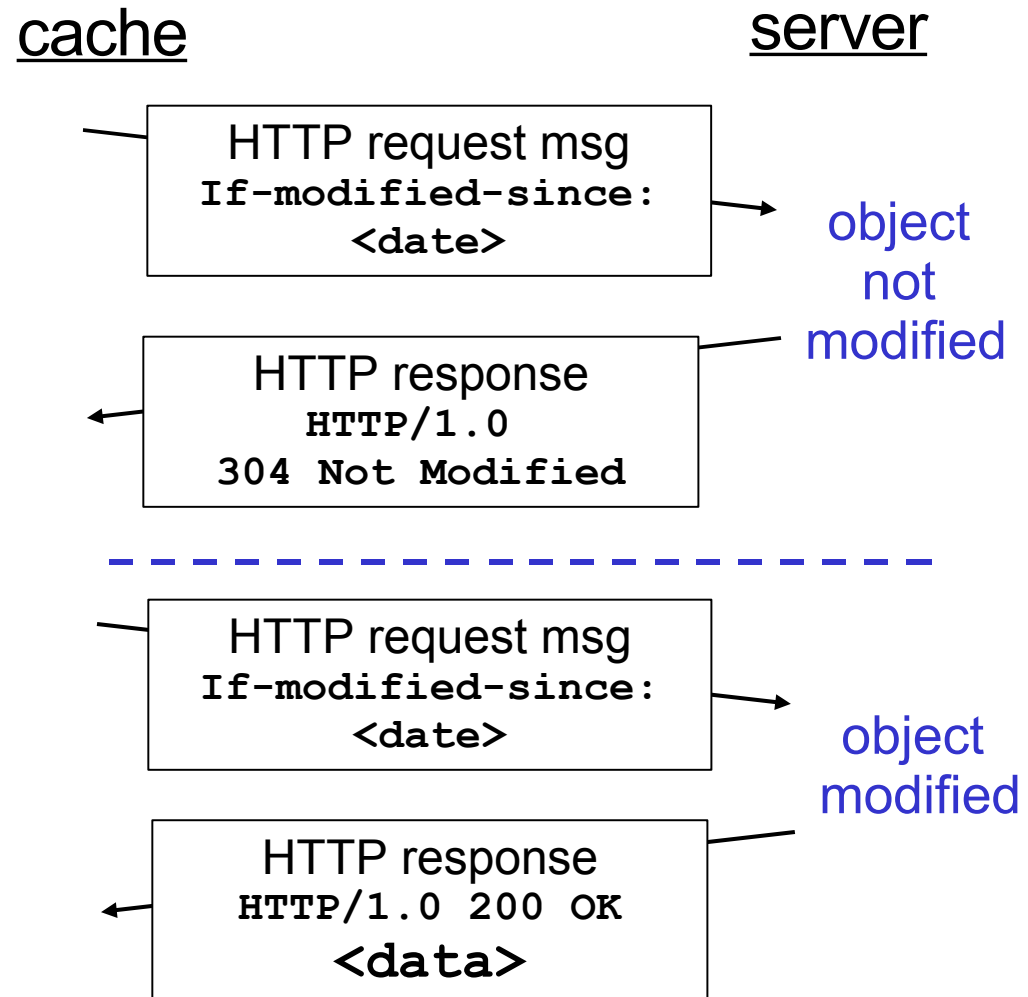
## Consequence

- 40% requests will be satisfied almost immediately
- 60% requests satisfied by origin server
- utilization of access link reduced to 60%, resulting in negligible delays (say 10 msec)
- total avg delay = Internet delay + access delay + LAN delay =  $.6 \cdot (2.01) \text{ secs} + .4 \cdot \text{milliseconds} < 1.4 \text{ secs}$



# Conditional GET

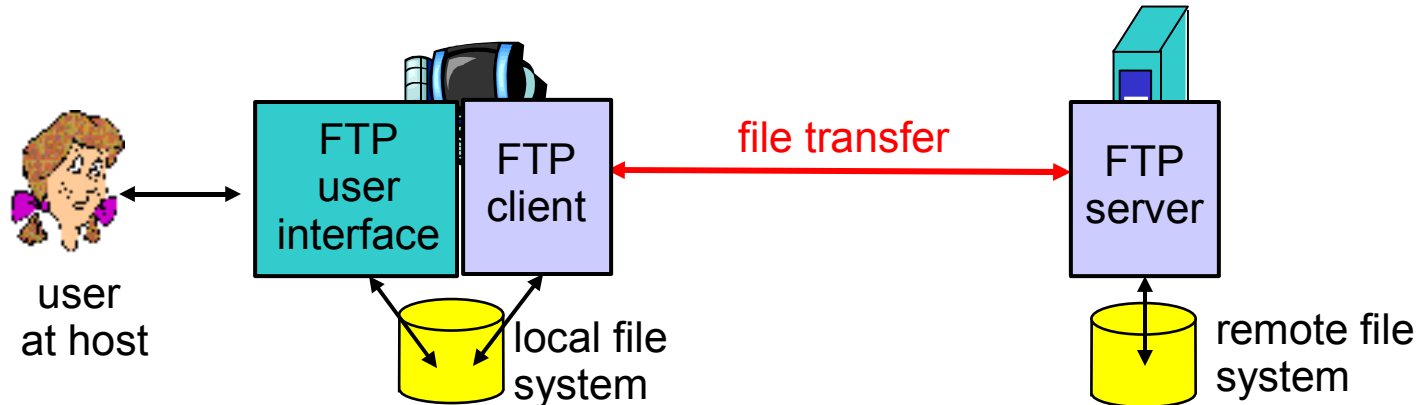
- **Goal:** don't send object if cache has up-to-date cached version
- cache: specify date of cached copy in HTTP request  
`If-modified-since:`  
`<date>`
- server: response contains no object if cached copy is up-to-date:  
`HTTP/1.0 304 Not Modified`



# Chapter 2: Application layer

- 2.1 Principles of network applications
- 2.2 Web and HTTP
- **2.3 FTP**
- 2.4 Electronic Mail
  - ◆ SMTP, POP3, IMAP
- 2.5 DNS
- 2.6 P2P file sharing
- 2.7 Socket programming with TCP
- 2.8 Socket programming with UDP
- 2.9 Building a Web server

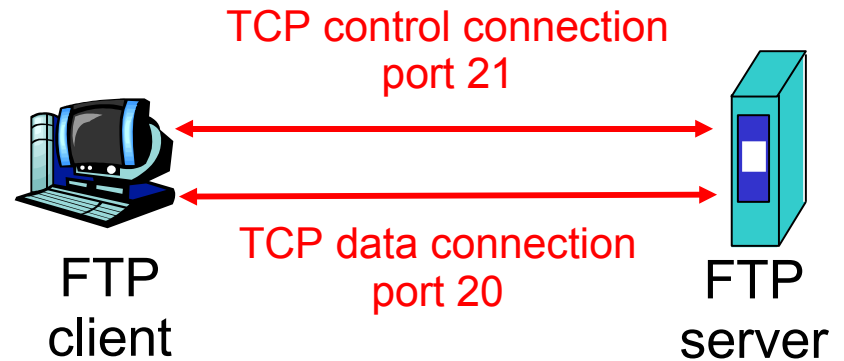
# FTP: the file transfer protocol



- transfer file to/from remote host
- client/server model
  - ◆ *client*: side that initiates transfer (either to/from remote)
  - ◆ *server*: remote host
- ftp: RFC 959
- ftp server: port 21

# FTP: separate control, data connections

- FTP client contacts FTP server at port 21, specifying TCP as transport protocol
- Client obtains authorization over control connection
- Client browses remote directory by sending commands over control connection.
- When server receives file transfer command, server opens 2<sup>nd</sup> TCP connection (for file) to client
- After transferring one file, server closes data connection.



- Server opens another TCP data connection to transfer another file.
- Control connection: “out of band”
- FTP server maintains “state”: current directory, earlier authentication



# FTP commands, responses

## Sample commands:

- sent as ASCII text over control channel
- **USER *username***
- **PASS *password***
- **LIST** return list of file in current directory
- **RETR *filename*** retrieves (gets) file
- **STOR *filename*** stores (puts) file onto remote host

## Sample return codes

- status code and phrase (as in HTTP)
- **331 Username OK, password required**
- **125 data connection already open; transfer starting**
- **425 Can't open data connection**
- **452 Error writing file**

# Recap

- Finish HTTP
- FTP

# Next time

- SMTP (email)
- DNS
- P2P