

# CS360 Theory of Computing

Alphabets, Strings, Languages,  
& Classes of Languages

# Alphabets & Strings

## Alphabet

- A finite set of symbols
- Usually denoted by  $\Sigma$

Example: a binary alphabet,  $\Sigma = \{0, 1\}$

## String (words)

- A sequence of alphabet symbols in an order

Example: 1010110

# String Operations

Length of a string:  $\mathbf{x} = |\mathbf{x}|$

- All our strings will have finite length

Empty string: denoted by  $\varepsilon$  (sometimes  $\Lambda$ )

- Length of empty string:  $|\varepsilon| = 0$

Power Op:

$\mathbf{x}^k = k$  copies of a string concatenated

Example:  $\mathbf{x} = 010$ ,  $\mathbf{x}^3 = 010010010$

Convention:  $\mathbf{x}^0 = \varepsilon$

# Prefixes, Suffixes & Substrings

Prefix: first  $k$  chars of a string  $v$  for some value  $k$

Suffix: last  $k$  chars of a string  $v$  for some value  $k$

Substring: a string  $u$  that is found as consecutive characters in a string  $v$

Concatenation:

Formal:  $x$  is a prefix of  $y$  if  $y = x \bullet z$  for some  $z$

- Often we will not write the “ $\bullet$ ”

# More String Operations

Counting members:

$n_b(\mathbf{x}) = \#$  of times alphabet symbol  $\mathbf{b}$  is found  
in string  $\mathbf{x}$

Reverse strings:

$\mathbf{rev}(\mathbf{x}) = \mathbf{x}$  written in reverse order (also,  $\mathbf{x}^r$ )

Palindrome:

$\mathbf{x}$  is a palindrome if  $\mathbf{x} = \mathbf{rev}(\mathbf{x})$

# Languages

$\Sigma^*$  is the set of all finite strings over alphabet  $\Sigma$

A language  $L$  is a set of strings over  $\Sigma$

- A subset of  $\Sigma^*$ , possibly  $\Sigma^*$  itself
- $\emptyset$  denotes the empty set over any  $\Sigma$

Note:  $\emptyset \neq \{\epsilon\}$

Complement:  $L' = \Sigma^* - L$

# Concatenation of Languages

If **L** and **M** are languages, then:

$$\mathbf{LM} = \{\mathbf{xy} \mid \mathbf{x} \in \mathbf{L} \text{ and } \mathbf{y} \in \mathbf{M}\}$$

Finite Power

$\mathbf{L}^k$  is  $k$  copies of **L** concatenated

Example:  $\mathbf{L}^3 = \mathbf{LLL}$

- $\mathbf{L}^0 = \{\varepsilon\}$
- $\mathbf{L}^1 = \mathbf{L}$

# Kleene Star

$$\mathbf{L}^* = \bigcup_{i=0}^{\infty} \mathbf{L}^i \quad \text{and} \quad \mathbf{L}^+ = \bigcup_{i=1}^{\infty} \mathbf{L}^i$$

If  $\mathbf{L} = \{a, b\}$  then:

$\mathbf{L}^* = \{\varepsilon, a, b, aa, ab, \dots\}$  i.e.  $\mathbf{L}^*$  includes  $\varepsilon$

$\mathbf{L}^+ = \{a, b, aa, ab, \dots\}$

Note:  $\mathbf{L}^{**} = \mathbf{L}^*$

$\mathbf{L}^*$  is **closed** under Kleene Star.

# Class of Languages

- A set of languages
- A set of sets (usually something in common)

Note: strings are members of languages, **not** members of classes of languages.

# Idea is to describe other languages

- Also have the normal set operations: union, intersection, difference, complement, etc.

Example:  $\mathbf{M} = (\mathbf{L} \cap \mathbf{N})^* \mathbf{OP}$

How do we test membership in  $\mathbf{M}$ ?

Each member of  $\mathbf{M}$  consists of

- A finite # of strings (0 or more) from  $\mathbf{L} \cap \mathbf{N}$  concatenated together
- Followed by a string from  $\mathbf{O}$
- Followed by a string from  $\mathbf{P}$