

CS 886: Game-theoretic methods for computer science

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Computer Science
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Outline

- 1 Introduction
 - Introduction
 - Two Communities
- 2 This Course
- 3 Examples
 - Selfish Routing
 - London Bus System

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Introduction

- Growth in settings where there are multiple self-interacting parties
 - Networks
 - Electronic marketplaces
 - Game playing
 - ...
- For participants to act optimally in such settings, they must take into account how other agents are going to act.
- We want to be able to
 - Understand the ways agents will interact and behave
 - Provide incentives so that agents behave the way **we** would like them to

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Two Communities

Economics

- Traditional emphasis on game theoretic rationality
- Describing how agents should behave
- Multiple self-interested agents

Computer Science

- Traditional emphasis on computational and informational constraints
- Building agents
- Individual or cooperative agents

New Research Problems

- How do we use game theory and mechanism design in computer science settings?
- How do we resolve conflicts between game-theoretic and computational constraints?
- Development of new theories and methodologies

New Research Area

Explosion of research in the area (Algorithmic game theory, computational mechanism design, Distributed algorithmic mechanism design, computational game theory,...)

- Papers appearing in AAI, AAMAS, UAI, NIPS, PODC, SIGCOMM, INFOCOMM, SODA, STOC, FOCS, ...
- Papers by CS researchers appearing in Games and Economic Behavior, Journal of Economic Theory, Econometrica,...
- Numerous workshops and meetings,...

This Course

The goals of this course

- Introduction to game theory and mechanism design
- Study how they are used in computer science
- Study computational issues that arise

Course structure

- Introductory lectures
- Current research papers

Topics

- Game theory
 - Normal form and extensive form games
 - Dominance and iterated dominance
 - Minimax strategies, Nash equilibrium, correlated equilibrium, backward induction, subgame perfect equilibrium, Bayesian games
 - Repeated games, Folk theorems
 - Coalitional game theory
- Social choice
 - Arrow's theorem, voting
- Mechanism Design
 - Incentive compatibility, individual rationality, positive and negative results, Revelation principle, Vickrey-Clarke-Groves mechanisms
 - Auctions (both single-item and combinatorial)
- Applications

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Prerequisites

- No formal prerequisites
- Students should be comfortable with mathematical proofs
- Some familiarity with probability
- Ideally students will have an AI course and background in algorithms and complexity
- **I will cover the game theory and mechanism design required**

Grading

- 2-3 assignments on game theory and mechanism design: 10%
- In class presentation(s): 20%
 - Peer-reviewed
- Class participation: 20%
- Research project: 50%

Presentations

Every student is responsible for presenting a research paper in class

- Short survey + a critique
- Everyone in class will provide feedback on the presentation
- Marks given on coverage of material + organization + presentation

Class Participation

You must participate!

- Before each class (before 6:00 am the day of the presentation) you must email me a list of comments on the paper to be presented¹
 - What is the main contribution?
 - Is it important? Why?
 - What assumptions are made?
 - What applications might arise from the results?
 - How can it be extended?
 - What was unclear?
 - ...

¹Plain text please. No attachments!

Projects

The goal of the project is to develop a deep understanding of a topic related to the course.

- The topic is open
 - Theoretical, experimental, *in-depth* literature review,...
 - Can be related to your own research
 - If you have trouble coming up with a topic, come and talk to me
- Proposals due October 18
- Final projects due December 6th²
- Students will present projects in class

²I will likely be flexible with this.

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Other Information

- Class times: Monday-Wednesday 10:00-11:30
- Office Hours: Mondays 1:30-2:30³
- Course website



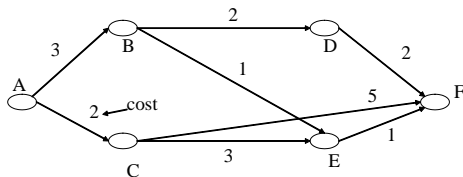
<http://www.cs.uwaterloo.ca/~klarson/teaching/F06->

³Except Sept 18, Oct 16, and Nov 20.

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Selfish Routing



- We want to find the least-cost route from S to T .
- Costs are private information – we do not know them
- We do know that agents (nodes) are interested in maximizing revenue
- How can we use this to figure out the least-cost route?

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London Bus System⁴

- 5 million passengers daily
- 7500 buses
- 700 routes
- The system has been privatized since 1997 by using competitive tendering
- *Idea*: Run an auction to allocate routes to companies

⁴As of April 2004

Auction Protocol

- Let G be set of all routes, I be the set of bidders
- Agent i submits bid $v_i(S)$ for all bundles $S \subseteq G$
- Compute allocation S^* to maximize sum of reported bids

$$V^*(I) = \max_{(S_1, \dots, S_n)} \sum_i v_i(S_i)$$

- Compute best allocation without each agent

$$V^*(I \setminus i) = \max_{(S_1, \dots, S_n)} \sum_{j \neq i} v_j^*(S_j)$$

- Allocate each agent S_i^* , each agent pays

$$P(i) = v_i^*(S_i^*) - [V^*(I) - V^*(I \setminus i)]$$

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London Bus System

- Mechanism: Generalized Vickrey Auction
 - Specifies the rules
 - Describes how outcome will be determined
- Strategies
 - Policies which specify what actions to take
 - Agents are self-interested and rational
- GVA is efficient and strategy-proof

Computational Issues

- Winner determination problem: Select bids to maximize sum of reported values
 - Maximum weighted set packing (NP-hard)
 - Solve this problem $I + 1$ times
- Agent valuation problem
- Communication complexity
 - Each agent has to communicate 2^{700} bids to the auctioneer

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